



How Facebook and Microsoft Successfully Leverage NVMe[™] Cloud Storage

Sponsored by NVM Express[™] organization, the owner of NVMe[™], NVMe-oF[™] and NVMe-MI[™] standards



Speakers

Ross Stenfort

Lee Prewitt











NVMe[™] In The Real World

Ross Stenfort, Hardware System Engineer

Facebook



facebook

Facebook's mission is to give people the power to build community and bring the world closer together.

Facebook @ Scale

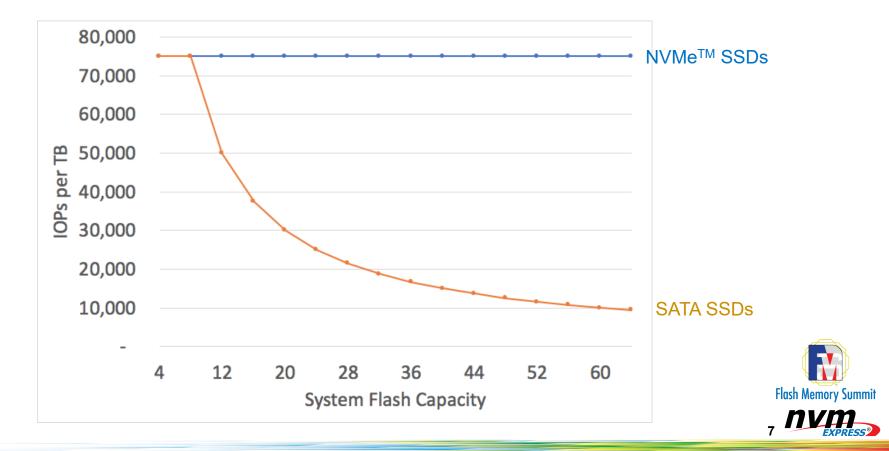
1 Billion

1.3 Billion





Hyperscale Requires IOPS to Scale with Capacity



NVMe[™] De-Allocate: Challenges and Improvements

- ➢ NVMe[™] De-allocate
 - Goal: It's a hint from the system to the SSD that the system is no longer tracking certain LBAs
 - Good
 - Reduces Write Ampliflication
 - Improves performance/endurance
 - Bad
 - Latency spikes due to De-allocate blocking Read/ Write
- Old Solution
 - Tune De-Allocate size on a system
 - Problem: The optimized de-allocate size varries based on supplier. Thus which supplier should I optimize for?
- Improved solution
 - > NVMe 1.4 allows the SSD to advertise it's preferred De-allocation size
 - If NSFEAT bit 4 = 0x1 then Namespace Perferred Deallocate Ganulativity (NPDG) is valid
 - > This allows systems to be optimized standard mechanisms.



Managing at Scale (1 of 3)

- > Challenge: Hyperscale Requires Debug with no physical access to the SSD.
- Challenge#1: Restricted access for vendor unique tools
- ➢ Solution:
 - NVMeTM CLI Open source with active industry contribution and updates
 - <u>https://github.com/linux-nvme/nvme-cli</u>
 - Vendor-unique CLI plugin that pulls and reports the logs in a common format

Challege#2: How do I get the debug information needed to resolve the issue Solution: Telemetry

- This allows SSD providers to get remote debug information to resolve issues
- Different data areas allows for different levels of debugging



Managing at Scale (2 of 3)

- Background: The amount of data written to a SSD may exceed the enduance of the SSD given the expected lifetime of the SSD. Given a fixed amount of write bandwidth a low the capacity SSD will wear out faster than a higher capacity SSD. Examples of applications where this can ocure are logging and caching.
- Challenge/ Real World Example:
 - Application only needs 256 GB but will use all the SSD capacity
 - Application write rate is high enough that it will wear out the 256 GB SSD
 - Application write rate scales per TB: Thus increasing capacity will not keep the SSD from wearing out.
- Solution: Namespace Management
 - Allows a 512 GB SSD to be configured as a 256GB SSD with double the endurance of a 256 GB SSD
 - Thus the application view is a 256GB with double the endurance



Managing at Scale (3 of 3)

- Challenge: How many blocks in my SSD have data and how many do not? If I de-allocate some blocks how many blocks really contain data? What is the effective over provisioning from a performance perspective?
- ➢ Solution:
 - Namespace Utilization (NUSE)
 - > Allows user to determine the number of LBAs that actually contain data.



Industry Challenge

- Security challenges are growing
 - NVM Express[™] supports SECURITY_SEND/ RECIEVE will allows for security protocols to be tunneled into NVM Express
 - There is even an open source tool for NVMe[™] Opal security:

https://github.com/Drive-Trust-Alliance/sedutil

• Secure Boot is also a common security requirement. This is a process that ensures the firmware running on the device is from the manufacture and not some other source.



Problem/ Industry call to action:

- There is no standard way to know if secure boot failed
- If firmware on a device is compromised, how is this identified vs any other type of failure?





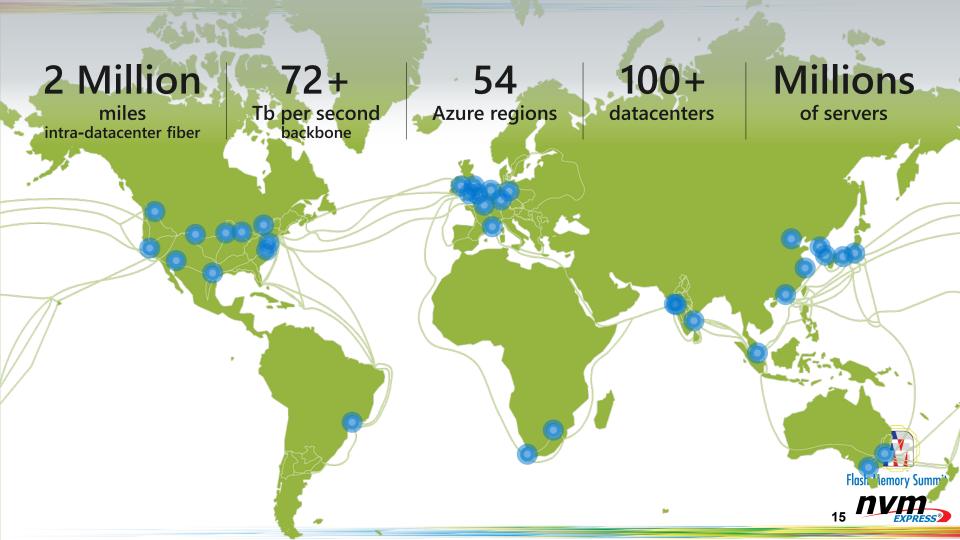


NVMe[™] at Hyper-Scale

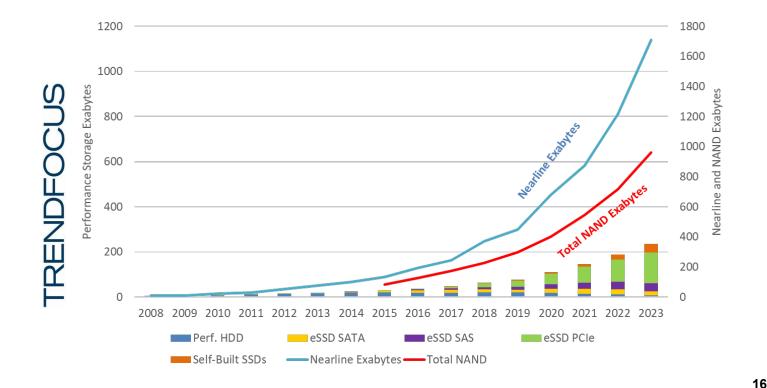
Lee Prewitt, Principle Hardware Program Manager Azure CSI - Microsoft



- Azure at a Glance
- Why NVMe[™]?
- Issues at Scale
 - Form factors
 - Need to allow for "rot in place"
 - Need for remote debugging
 - Need for security



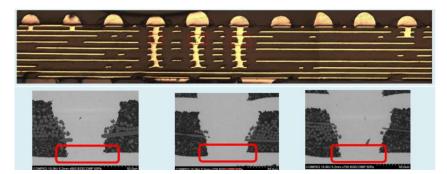
Why NVMe[™]? - Exploding Storage Growth





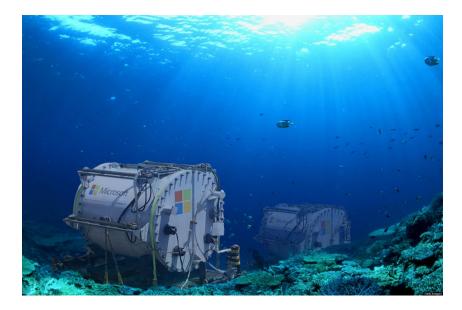
Issues at Scale – Form Factors

- m.2 has run its course
 - Power and thermal constraints
 - Fragile PCB and connector
 - Not hot-swappable
- E1.L and E1.S are here to replace it
 - Built from the ground up for datacenter use cases
- Good news is that they support NVMe[™] too!





Issues at Scale – Need to allow for "Rot in Place"



Use the Endurance and Performance metrics for auto tiering

- Allows for fitting the workload to the device
- Allows for the ability to adjust the temperature of the data over time
- Allow for 5 7 year device service life

Zoned Name Spaces for QLC

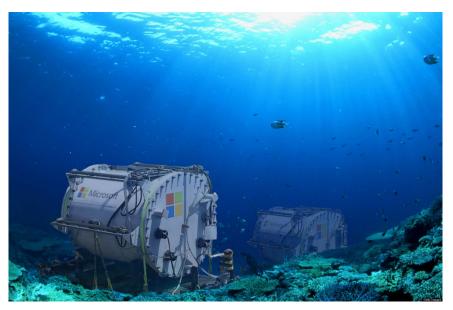
- Reduce WAF due to large sequential writes
- Reduce DRAM due to large indirection unit
- Reduce overprovisioning due to minimal garbage collection



Issues at Scale – Need for Remote Debugging

- Timestamp
 - Drive events correlated to system (BIOS and OS) events
- Telemetry
 - Host initiated IO failures
 - Drive Initiated Firmware panic?
- SMART
 - Both standard and vendor unique collected once an hour
 - Hey SSD IHVs. How many terabytes would you like to see?

Caveat: Any data that leaves the datacenter must be in human readable form!





Issues at Scale – Need for Security



eDrive on Windows

Opal v2 plus IEEE 1667 secure silo

Hardware Root of Trust

- Secure boot
- Signed firmware
- Cerberus

Device Hardening

- Pen and Fuzz testing
- Locking of debug ports and vendor unique commands



Questions?







Architected for Performance

